

Please provide complete and well-written solutions to the following exercises.

Due October 25, in the discussion section.

Homework 4

Exercise 1. Let A be a real $m \times n$ matrix. Let $x \in \mathbf{R}^n$ and let $b \in \mathbf{R}^m$. Show that the function $f: \mathbf{R}^n \rightarrow \mathbf{R}$ defined by $f(x) = \frac{1}{2} \|Ax - b\|^2$ is convex. Moreover, show that

$$\nabla f(x) = A^T(Ax - b), \quad D^2f(x) = A^T A.$$

Exercise 2. Let A be an $m \times n$ real matrix with $m \geq n$. Then A has rank n if and only if $A^T A$ is positive definite.

(Hint: $A^T A$ is automatically positive semidefinite by a previous exercise.)

Exercise 3. Show the following identity. Let A be an $r \times r$ real matrix, let U be an $r \times s$ real matrix, and let V be an $s \times r$ real matrix. Assume that A is invertible and that $I + VA^{-1}U$ is invertible, where I is the $s \times s$ identity matrix. Then $A + UV$ is invertible and

$$(A + UV)^{-1} = A^{-1} - (A^{-1}U)(I + VA^{-1}U)^{-1}(VA^{-1}).$$

In particular, if $s = 1$, we get the Sherman-Morrison formula:

$$(A + UV)^{-1} = A^{-1} - \frac{A^{-1}UV A^{-1}}{1 + VA^{-1}U}.$$

Exercise 4. Give a bound for the number of arithmetic operations needed in Algorithm ???. Assume that P_n is known, so that computing P_n does not require any arithmetic operations. (Hint: your bound should be something like $(m - n)n^2$.) Compare this bound to simply minimizing $\|Ax - b\|^2$ directly. (In that case, you should need about mn^2 arithmetic operations.)

The key point here is that, once P_n is known, recursive least squares is much better. For example, if $m - n = 10$, then recursive least squares requires around $10n^2$ arithmetic operations. But minimizing $\|Ax - b\|^2$ directly would require mn^2 operations, which is much larger.

Exercise 5. Let $x^{(1)}, \dots, x^{(m)}$ be vectors in \mathbf{R}^2 such that $\|x^{(i)}\| = 1$ for all $1 \leq i \leq m$. Show that there exist $i, j \in \{1, \dots, m\}$ with $i \neq j$ such that $\langle x^{(i)}, x^{(j)} \rangle > 1 - 100/m^2$.

Exercise 6 (Logistic Regression). Let $x^{(1)}, \dots, x^{(m)} \in \mathbf{R}^n$ and let $y_1, \dots, y_m \in \{0, 1\}$. For the sake of intuition, we can think of each vector $x^{(i)}$ as a vector of words in an email, and y_i classifies email $i \in \{1, \dots, m\}$ as either spam ($y_i = 1$) or not spam ($y_i = 0$). Given this data, we would like to find a way to classify future emails as spam or not spam. (This is what a spam filter does.) For any $t \in \mathbf{R}$, define the logistic function $g: \mathbf{R} \rightarrow (0, 1)$ by

$$g(t) = \frac{1}{1 + e^{-t}}.$$

The function g is meant to be a differentiable approximation to a function whose output is either 0 or 1.

First, verify that $g'(t) = g(t)(1 - g(t))$ for any $t \in \mathbf{R}$. Then, consider the log-likelihood function $L: \mathbf{R}^n \rightarrow \mathbf{R}$ defined by

$$L(z) := \log \left(\prod_{i=1}^m [g(\langle z, x^{(i)} \rangle)]^{y_i} [1 - g(\langle z, x^{(i)} \rangle)]^{1-y_i} \right), \quad \forall z \in \mathbf{R}^n.$$

We would like to maximize L . The idea here is that if L is large, then z is a set of parameters (or “weights”) that accurately classifies known emails as spam or not spam. So, once we find z , and if we have some new email $x \in \mathbf{R}^n$, then $g(\langle z, x^{(i)} \rangle) \approx 0$ means the new email is probably not spam, and $g(\langle z, x^{(i)} \rangle) \approx 1$ means the new email is probably spam.

Show that

$$\nabla L(z) = \sum_{i=1}^m (y_i - g(\langle z, x^{(i)} \rangle)) x^{(i)}, \quad \forall z \in \mathbf{R}^n.$$

This computation then gives the formula for a gradient ascent method for maximizing L .

Exercise 7. Define

$$A = \begin{pmatrix} 1 & 2 \\ 3 & 4 \\ 1 & 3 \\ 2 & 3 \end{pmatrix}, \quad b = \begin{pmatrix} 1 \\ 3 \\ 4 \\ 6 \end{pmatrix}.$$

Minimize the function $\|Ax - b\|^2$ over all $x \in \mathbf{R}^2$, either by hand, or using a computer program that you write yourself. In either case, use the recursive least squares method.

Verify that the x you found does actually minimize $\|Ax - b\|^2$.

Exercise 8. Let $c, x \in \mathbf{R}^n$, let $b \in \mathbf{R}^m$, and let A be an $m \times n$ real matrix.

It is possible to essentially put all of the variables of a linear program into the constraint. Show that the following linear program is equivalent to the standard one. (That is, the minimum/maximum values and the x achieving the minimum/maximum value is the same for both linear programs.)

$$\begin{aligned} & \text{maximize } t \quad \text{subject to the constraints} \\ & \{t \in \mathbf{R}: t \leq c^T x, \forall x \in \mathbf{R}^n \text{ such that } Ax = b, x \geq 0\}. \end{aligned}$$

Exercise 9. Let $c, x \in \mathbf{R}^n$, let $b \in \mathbf{R}^m$, and let A be an $m \times n$ real matrix. Show that the linear program

$$\begin{aligned} & \text{minimize } c^T x \quad \text{subject to the constraints} \\ & Ax \leq b. \end{aligned}$$

is equivalent to the following linear program in standard form

$$\text{minimize } \begin{pmatrix} c \\ -c \\ 0 \end{pmatrix}^T \begin{pmatrix} x^+ \\ x^- \\ z \end{pmatrix} \quad \text{subject to the constraints}$$

$$(A \quad -A \quad I) \begin{pmatrix} x^+ \\ x^- \\ z \end{pmatrix} = b, \quad \begin{pmatrix} x^+ \\ x^- \\ z \end{pmatrix} \geq 0.$$

By equivalent, we mean that if $x \in \mathbf{R}^n$, and if we define $x^+ := \max(x, 0)$, $x^- := \max(-x, 0)$, then $x = x^+ - x^-$. Similarly, if $x^+, x^- \geq 0$, then we define $x = x^+ - x^-$. And the minimum values and the x achieving the minimum value for both linear programs is the same. (Here the maximum is defined component-wise, e.g. if $x = (-1, 2, 3)$, then $\max(x, 0) = (0, 2, 3)$.)